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- (S) References cited:
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  US-A- 4 776 028
  HELD G. 'DATA COMPRESSION TECHNIQUES AND APPLICATIONS HARDWARE AND SOFTWARE CONSIDERATIONS' 1983 , J. WILEY AND SONS , CHICESTER, GB

- (3) Proprietor: Hewlett-Packard Company Mail Stop 20 B-O, 3000 Hanover Street Palo Alto, California 94304 (US)
- (72) Inventor: Miller, Steven O. 14818 NE 78th Street Vancouver Washington 98682 (US)
- (4) Representative: Liesegang, Roland, Dr.-Ing. et al FORRESTER & BOEHMERT Franz-Joseph-Strasse 38 D-80801 München (DE)

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Description

## Technical Field

This invention relates to a method of graphics data compression and, more particularly, to a method capable of compressing, selectively, data on a raster-row-by-raster-row basis or on a byte-group-by-byte-group basis.

## **Background Art**

Bytes of data often describe graphics images as an arrangement of raster rows. A raster row is one line or strip of a graphics image. Numerous parallel raster rows may be arranged or stacked to produce an entire graphics image.

Each byte of data describes a part of a raster row within the graphics image. Because each byte describes a part of a raster row, the data itself may be referred to as a raster row of data, or data arranged on a raster-row-by-raster-row basis, as is done in the following text.

Because graphics images are often complex, the data describing the images may be extensive and detailed. Accordingly, transferring such data between components in a system, such as between a host computer and a printer, often requires a large amount of memory and may result in an undesirable time lag. To limit the required memory and to prevent such time lags, the data may be compressed so that information is transferred by employing less data, with time thereby saved.

Data may be compressed in a variety of ways. For example, Gilbert Held, in his book <u>Data Compression</u>, published by Wiley Hayden, copyright 1984, discloses a compression method wherein bits of data within one raster row are compared with bits of data in an adjacent raster row. By comparing the bits, the differences between the rows may be noted. Once noted, one raster row may be described by stating how the bits in that row differ from the bits in the other row. The data describing the differences may be less extensive than the data describing the raster row, resulting in a compression of data.

However, comparing different raster rows of data on such a bit-by-bit basis does not recognize repetitious data within the raster rows themselves. Additionally, such a comparison does not lend itself to a concise method of recording the bit changes between the rows.

From US-A- 4 724 483 or US-A- 4 776 028 it is known to subdivide binary data into units of e.g. 16 bits each for compression on a unit-to-unit basis and on a row-to-row basis.

It is the problem of the present invention to provide a data compression method which enables recognizing repetitions within raster rows as well as recording changes between rows concisely.

The invented graphics images data compression method defined in claim 1 essentially selectively compresses data on a raster-row-by-raster-row basis or on a byte-group-by-byte-group basis. The raster-row-by-raster-row basis compares bytes within adjacent raster rows and thereby detects repetitions between the rows. When data is compared on a raster-row-by-raster-row basis, the invented method records any differences as relative offset and replacement bytes. The byte-group-by-byte-group basis compares different byte groups within a single raster row, and thereby detects repetitions within the row itself. Selectively compressing data according to such bases results in superior data compression.

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## **Brief Description of the Drawings**

Figure 1 is a block diagram showing the invented raster-row-by-raster-row compression method.

Figure 2 depicts the relationship of Figures 3A, 3B, 3C.

Figures 3A-C constitute a detailed block diagram showing the invented data compression method.

# <u>Detailed Description and Best Mode for Carrying</u> Out the Invention

Figure 1 shows generally at 10 a block diagram of the invented raster row data compression method. Initially, a block of data is arranged on a raster-row-by-raster-row basis, as signified by block 12. Again, a block of data arranged on a raster-row-by-raster-row basis refers to data which describes an image as an arrangement of raster rows.

Step 14 then compares a first and second raster row of data and the changes between the rows are noted at step 16. Specifically, each byte in the second raster row that differs from the corresponding byte in the first raster row is recorded, at step 18, as relative offset and replacement bytes. Corresponding bytes are bytes of data that define parts of an image which are vertically adjacent but in different raster rows. Relative offset and replacement bytes refer to a manner of recording the changes between such raster rows.

For example, the number of bytes to be replaced in a second raster row because they differ from the corresponding bytes in the first raster row, may be recorded as a three-bit number. The location of the bytes to be replaced may be recorded as a five-bit number. The five-bit number is the relative offset value. The offset is the number of bytes in the present raster row from the last untreated byte, or in other words, the number of bytes from the byte following the last replaced byte. Alternatively, the offset may be relative to the left graphics margin.

An offset that is large may require several bytes

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to represent. The numbers zero through thirty may signify a relative offset of zero to thirty where zero refers to the next uncompressed byte. If the five-bit number is thirty-one, then the next byte is interpreted as an additional offset value and is added to the previous five-bit number. For example, an offset of thirty-one may be recorded as the number thirty-one plus an additional offset of zero contained in the next byte of data. Thirty-two may be recorded as the number thirty-one plus an additional offset of one. If such next byte contains the number two hundred fifty-five, then such value is added to the five-bit number and the next byte is included as an offset, and so on.

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As explained, the number of bytes to replace and the location (offset) of such bytes are recorded. The replacement bytes themselves follow the bytes defining the offset value. For each byte that is replaced, a replacement byte must follow. The number of replacement bytes may range from one to eight, but may not exceed eight for a single command byte because only three binary digits are available to indicate the number of bytes to replace. Accordingly, the format of such data may be expressed as:

# <command byte> [<optional offset bytes>]<replacement bytes>

where the command byte includes the three-bit number of bytes to replace and the initial five-bit offset value. For example, in binary code, if the command byte is "0100 0000", then the next three bytes in the present raster row are to be replaced with whatever replacement bytes follow. If the command byte is "0110 0101", then the next five bytes in the present raster row are jumped over and the following four bytes are replaced with the appropriate four replacement bytes. If the command byte is "0000 0000", the next byte is to be replaced.

In the preferred embodiment, compiling changes between raster rows or byte groups refers to recording the number and location of such changes as explained. However, different recording methods may be used without departing from the invented compression method.

Data that does not change between the compared raster rows is replicated at step 20. The replicated bytes are not part of the compressed data because they would already exist in a task processor's memory or buffer. In other words, the step of replicating simply leaves data unchanged. The replicated bytes and the relative offset and replacement bytes completely describe the data within the second raster row. The data may then be sent to a task processor as symbolized by step 22. Again, only the offset and replacement bytes are sent to the task processor because the replicated data is already there and remains unchanged. Step 24 then determines whether the entire block of data has been compressed. If so, step 26 ends the compression. If not, the method recommences at step 14 by comparing subsequent raster rows of information.

Initially, a first and second raster row were compared. Comparing data describing subsequent raster rows means that what was previously referred to as the "second" raster row is now the "first" raster row. In other words, "first" raster row, as used herein, means the reference or seed row. A next raster row becomes the second raster row, and is compared to the uncompressed data which describes the new first raster row. Accordingly, the terms first and second raster rows refer to different rows as the compression progresses.

Figure 2 depicts the relationship between Figures 3A-C, wherein Figure 3A is depicted by block 28, Figure 3B by block 30 and Figure 3C by block 32. Figures 3A, 3B and 3C show a block diagram of the invented selectable compression method.

In Figure 3A, block 34 represents bytes of data describing raster rows of an image. Step 36 identifies the first raster row of such data.

The first raster row at the beginning of compression activity is then compared with a reference raster row as shown at step 38. The reference raster row is generated by the system employing the invented method, and would typically be set to a value of zero because the initial rows of a graphics image are generally blank.

The next step 40 is to note and compile the changes between the first raster row and the reference raster row. The compilation of such changes constitutes a compressed description of the data describing the first raster row. Step 42 then determines whether the compression was sufficient. A compression is sufficient if it meets a preselected compression ratio, for example 9:1.

If the compression is not sufficient, then the invented method determines, at step 44, whether the bytes of data within the first raster row have been compressed on a byte-group-by-byte-group basis. If a byte-group compression has not been done, then the method compares byte groups within the first raster row, as shown at step 46. A byte group is a preselected number of bytes within a given raster row. Steps 40 and 42 are then repeated for the byte-group compression.

If the raster row compression was insufficient, and the byte-group compression was also insufficient, then the invented method uses the most compressed data, as shown by step 48. The most compressed data may be the raster row compression, the byte-group compression or the uncompressed data itself, whichever requires less memory to store.

If either the initial raster row compression was sufficient, the subsequent byte-group compression was sufficient or if the most compressed data has been selected for use, the next step 50 is to convey any such data to the task processor. The invented method then determines at step 52 whether the com-

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pression is complete. If so, the task is ended at step 54. If not, the method proceeds with the steps labeled "A" at 56.

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Figure 3B depicts the steps following the label "A". Initially, the method identifies the next raster row at step 58. That row is then compressed on the same basis as was the previous raster row. For example, if the previous raster row was compressed on a bytegroup-by-byte-group basis, the next raster row will also be compressed on a byte-group-by-byte-group basis. Step 60 determines what the previous basis of compression was. If the previous raster row was uncompressed, then either compression method may be selected.

If the previous compression was on a raster-rowby-raster-row basis, the invented method proceeds through the steps labeled "B" at 62. If the previous compression was on a byte-group-by-byte-group basis, the invented method proceeds with the steps labeled "C" at 64.

The steps identified by label "B" begin by comparing the present raster row with the previous raster row, as shown by step 66. The changes between the rows are noted and compiled at step 68. The method then determines, at step 70, whether the compression was sufficient. If so, the data resulting from the raster row compression is used as shown in Figure 3C at step 72.

If step 70 determines that there was not a sufficient compression, then the next step is to determine whether the raster row has been compressed on a byte-group-by-byte-group basis, as shown at 74. If not, the method proceeds to those steps following label "C" at 64.

Accordingly, the steps following label "C" will be performed if the previous raster row was compressed on a byte-group-by-byte-group basis, as determined in step 60, or if the raster row compression of steps 66, 68 did not result in a sufficient compression. In either case, byte groups are compared at 76. The changes between byte groups are noted and compiled at 78.

The method then determines, at step 80, whether the byte-group compression was sufficient. If so, the data compressed on the byte-group-by-byte-group basis will be used, as shown in Figure 3C at step 82. If the byte-group compression did not result in a sufficient compression, the method determines whether the data has been compressed on a raster-row-byraster-row basis, as shown in step 84. If not, then the method continues with the steps following label "B" at

Figure 3C further depicts the invented data compression method. If a raster row has been compressed on a raster-row-by-raster-row basis and a byte-group-by-byte-group basis, and neither compression resulted in a sufficient compression, then the method will advance to block 86 and use the most compressed data. The most compressed data may be either the raster row compressed data, the byte-group compressed data or the uncompressed data.

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Whichever data is used, it is conveyed to task processor at step 88. The method then determines, at step 90, whether the data compression is complete, and if so the task is ended, as shown in block 92. If the compression is not complete, the method returns to label "A" at 56 in Figure 3B and proceeds with the subsequent steps.

## Industrial Applicability

The raster row data compression method and the selectable data compression method are applicable to any data describing graphics images on a rasterrow-by-raster-row basis. Compressing data by comparing raster rows has a high compression ratio on data that is repetitive from raster row to raster row, or in the "vertical" direction. Byte-group data compression has a high compression ratio for data that is repetitive within a single raster row, or in the "horizontal"

The selectable data compression method allows for data to be compressed on either a raster row or a byte-group basis, depending on which method of compression produces sufficient results. Such a selective compression method results in superior performance compared to other data compression methods because graphics images are often repetitive in a horizontal or vertical direction, but not always in both directions at the same time.

## Claims

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1. A method of graphics data compression used in a system, wherein byte groups of data are arranged on a raster-row-by-raster-row basis, for compressing such data comprising:

comparing (66,76) such raster rows of data selectively either on a raster-row-by-rasterrow adjacency basis, or on a byte-group-by-bytegroup adjacency basis;

utilizing a selected one of such comparison bases, and noting (68,78) the changes between the compared adjacent rows or adjacent byte groups;

compiling (68,78) such changes so that they constitute a first compressed description of one of the compared rows or byte groups;

determining (70,80) whether such first compressed description results in a sufficient compression of the data;

if so, using (72,82) the first compressed description:

if not, comparing (66,76) the data on the other comparison basis not used in the initial

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comparison;

noting (68,78) the changes between the compared adjacent rows or adjacent byte groups; compiling (68,78) such changes so that they constitute a second compressed description of one of the compared rows or byte groups;

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determining (70,80) whether such second compressed description results in a sufficient compression of the data;

if so, using (72,82) the second compressed description;

if not, using (86) the most compressed description between the first and second compressed descriptions and the uncompressed data; and

repeating the above-disclosed steps for subsequent raster rows of data.

- 2. The method of claim 1, wherein the step of comparing (76), when it is done on a byte-group-by-byte-group basis, comprises comparing a predetermined number of bytes of data with a predetermined number of successively adjacent bytes of data within the same raster row, and repeating such comparison of bytes until all the bytes comprising the raster row have been compared.
- The method of claim 1 or 2, wherein the step of repeating comprises comparing subsequent raster rows of data in the same method as the immediately prior data was compared, i.e., a rasterrow-by-raster-row basis or a byte-group-by-bytegroup basis depending on how the immediately previous data was compared.
- 4. The method of one of claims 1 to 3 further comprising beginning the data compression by comparing the data on a raster-row-by-raster-row basis, performing the steps of noting, compiling and determining, and then, if not sufficiently compressed, comparing the data on a byte-group-by-byte-group basis for a predetermined number of raster rows, performing the steps of noting, compiling, determining and using, and repeating the above-disclosed steps by recommencing with the raster-row-by-raster-row comparison.
- 5. The method of one of claims 1 to 4, wherein the steps of comparing (66), noting (68) and compiling (68), when they are done on a raster-row-by-raster-row basis, further comprise recording each byte in a second raster row that differs from the corresponding byte in a first raster row as relative offset and replacement bytes, and compiling such offset and replacement bytes so that they constitute a compressed description of such second raster row.

Patentansprüche

 Verfahren zur Komprimierung von Graphikdaten, welche in einem System verwendet werden, wobei Bytegruppen von Daten rasterzeilenweise angeordnet werden, um diese Daten zu komprimieren, mit den Verfahrensschritten:

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Vergleichen (66, 76) solcher Rasterzeilen von Daten wahlweise auf der Basis benachbarter Rasterzeilen oder auf einer der Basis benachbarter Bytegruppen,

Verwenden einer ausgewählten dieser Vergleichsbasen und Notieren (68, 78) der Veränderungen zwischen verglichenen benachbarten Zeilen oder benachbarten Bytegruppen,

Kompilieren (68, 78) dieser Veränderungen, so daß sie eine erste komprimierte Beschreibung einer der verglichenen Zeilen oder Bytegruppen bilden,

Bestimmen (70, 80), ob diese erste komprimierte Beschreibung eine ausreichende Komprimierung der Daten ergibt, wenn ja, Verwenden (72, 82) der ersten komprimierten Beschreibung,

wenn nein, Vergleichen (66, 76) der Daten auf der anderen Vergleichsbasis, welche beim anfänglichen Vergleich nicht verwendet wurde,

Notieren (68, 78) der Veränderungen zwischen den verglichenen benachbarten Zeilen oder benachbarten Bytegruppen, Kompilieren (68, 78) dieser Veränderungen, so daß sie eine zweite komprimierte Beschreibung einer der verglichenen Zeilen- oder Bytegruppen bilden,

Bestimmen (70, 80), ob diese zweite komprimierte Beschreibung eine ausreichende Komprimierung der Daten ergibt,

wenn ja, Verwenden (72, 82) der zweiten komprimierten Beschreibung,

wenn nein, Verwenden (86) der von der ersten und der zweiten komprimierten Beschreibung und den nicht komprimierten Daten am meisten komprimierten Beschreibung, und Wiederholen der obenbeschriebenen Schritte für nachfolgende Rasterzeilen von Daten.

- Verfahren nach Anspruch 1, bei dem der Schritt
  des Vergleichens (76) dann, wenn er bytegruppenweise ausgeführt wird, das Vergleichen einer
  vorgegebenen Anzahl von Datenbytes mit einer
  vorgegebenen Anzahl von nachfolgenden benachbarten Datenbytes innerhalb derselben Rasterzeile umfaßt, und dieser Bytevergleich wiederholt wird, bis alle Bytes in der Rasterzeile verglichen worden sind.
  - Verfahren nach Anspruch 1 oder 2, bei dem der Schritt des Wiederholens das Vergleichen nachfolgender Rasterzeilen von Daten gemäß denselben Verfahren wie beim unmittelbar vorherge-

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henden Datenvergleich umfaßt, d.h., rasterzeilenweise oder bytegruppenweise abhängig davon, wie die unmittelbar vorhergehenden Daten verglichen worden sind.

- 4. Verfahren nach einem der Ansprüche 1 bis 3, bei dem ferner die Datenkomprimierung begonnen wird, indem die Daten rasterzeilenweise verglichen werden, die Schritte Notieren, Kompilieren und Bestimmen durchgeführt werden und dann, wenn die Komprimierung noch nicht ausreichend ist, die Daten bytegruppenweise für eine vorgegebene Anzahl von Rasterzeilen verglichen werden, die Schritte Notieren, Kompilieren, Bestimmen und Verwenden durchgeführt werden und die obenbeschriebenen Schritte wiederholt werden, indem mit dem Rasterzeilen-Vergleich wiederbegonnen wird.
- 5. Verfahren nach einem der Ansprüche 1 bis 4, bei dem die Schritte des Vergleichens (66), Notierens (68) und Kompilieren (68) dann, wenn sie rasterzeilenweise ausgeführt werden, ferner das Aufzeichnen jedes Bytes in einer zweiten Rasterzeile, welches von dem entsprechenden Byte in einer ersten Rasterzeile abweicht, als relative Offset- und Ersatzbytes und das Kompilieren dieser Offset- und Ersatzbytes umfaßt, so daß diese eine komprimierte Beschreibung der zweiten Rasterzeile bilden.

#### Revendications

 Procédé de compression de données graphiques utilisé dans un système, dans lequel des groupes d'octets de données sont disposés sur une base de rangée de trame par rangée de trame, afin de compresser ces données, comprenant les étapes consistant:

à comparer (66, 76) ces rangées de trame de façon sélective soit sur une base contiguë de rangée de trame par rangée de trame, ou sur une base contiguē de groupe d'octets par groupe d'octets;

à utiliser une base sélectionnée parmi ces bases de comparaison, et à noter (68, 78) les changements entre les rangées adjacentes comparées ou les groupes d'octets adjacents comparés;

à compiler (68, 78) ces changements de façon à ce qu'ils constituent une première description compressée de l'une des rangées comparées ou de l'un des groupes d'octets comparés;

à déterminer (70, 80) si oui ou non cette première description compressée aboutit à une compression suffisante des données; si c'est oui, à utiliser (72, 82) la première description compressée;

si c'est non, à comparer (66, 76) les données sur l'autre base de comparaison non utilisée dans la comparaison initiale;

à noter (68, 78) les changements entre les rangées adjacentes comparées ou les groupes d'octets adjacents comparés;

à compiler (68, 78) ces changements de façon à ce qu'ils constituent une seconde description compressée de l'une des rangées comparées ou de l'un des groupes d'octets comparés;

à déterminer (70, 80) si oui ou non cette seconde description compressée aboutit à une compression suffisante des données;

si c'est oui, à utiliser (72, 82) la seconde description compressée;

si c'est non, à utiliser (86) la description la plus compressée parmi les première et seconde descriptions compressées et les données non-compressées; et

à répéter les étapes divulguées ci-dessus pour des rangées de trame ultérieures.

- 2. Procédé selon la revendication 1, dans lequel l'étape de comparaison (76), lorsqu'elle est effectuée sur une base de groupe d'octets par groupe d'octets, comprend une comparaison d'un nombre prédéterminé d'octets de données avec un nombre prédéterminé d'octets successivement adjacents de données dans la même rangée de trame, et une répétition de cette comparaison d'octets jusqu'à ce que tous les octets constituant la rangée de trame aient été comparés.
- 3. Procédé selon la revendication 1 ou 2, dans lequel l'étape de répétition comprend la comparaison de rangées de trame ultérieures selon le même procédé que celui au moyen duquel les données immédiatement précédentes ont été comparées, c'est-à-dire, sur une base de rangée de trame par rangée de trame ou sur une base de groupe d'octets par groupe d'octets en fonction de la façon selon laquelle les données précédentes ont été comparées.
- 4. Procédé selon l'une des revendications 1 à 3 comprenant également le démarrage de la compression de données par une comparaison des données sur une base de rangée de trame par rangée de trame, la réalisation des étapes consistant à noter, à compiler et à déterminant, et ensuite, si elles ne sont pas suffisamment compressées, la comparaison des données sur une base de groupe d'octets par groupe d'octets pour un nombre prédéterminé de rangées de trame, la réalisation des étapes consistant à noter,

à compiler, à déterminer et à utiliser, et la répétition des étapes divulguées ci-dessus en recommençant avec la comparaison de rangée de trame par rangée de trame.

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5. Procédé selon l'une quelconque des revendications 1 à 4, dans lequel les étapes consistant à comparer (66), à noter (68) et à compiler (68), lorsqu'elles sont réalisées sur une base de rangée de trame par rangée de trame, comprennent également un enregistrement de chaque octet dans une seconde rangée de trame qui diffère de l'octet correspondant dans une première rangée de trame sous la forme d'un décalage relatif et d'octets de remplacement, et une compilation de ce décalage et de ces octets de remplacement de façon à ce qu'ils constituent une description compressée de cette seconde rangée de trame.

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